Reference fallible endgame play

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A reference model of fallible endgame play is defined in terms of a spectrum of endgame players whose play ranges in competence from random to optimal choice of move. They may be used as suitable practice partners, to differentiate between otherwise equi-optimal moves, to promote or expedite a result, to assess an opponent, to run Monte Carlo simulations, and to identify the difficulty of a position or a whole endgame.

1. INTRODUCTION

When a player is armed with perfect information from an endgame table (EGT), questions still arise as to how to play a fallible opponent. In a drawn position, how best can the attacker pressure the defender or the defender resist the attacker (Levy 1987, 1991; Nunn, 2002; Schaeffer 1991, 1997)? How can the opponent be given the best opportunity to cede depth in a situation where a draw claim is likely (Haworth, 2000, 2001)? How can a player identify and adapt to apparent opponent fallibility? More generally, one might ask how well a fallible player is likely to perform in an endgame, and how ‘difficult’ that endgame or a specific position is.

A model of fallible endgame play is defined here, providing a way to address these questions. It defines a spectrum of reference endgame players, REPs, with defined levels of competence ranging from random to optimal play. A ‘better’ move (in terms of value and depth) is at least as likely to be chosen as a ‘worse’ move.

Section 2 introduces the key concepts while section 3 states the first results. Section 4 demonstrates some basic uses of fallibility and section 5 demonstrates how to assess the competence of the opponent by revisiting the historic Browne-BELLE KQKR games. Section 6 surveys ways of exploiting these reference fallible players and a fallible opponent.

2. KEY CONCEPTS

White here will, as is customary, be the winning or attacking player: Black is losing or defending. In scenario BFL, Black is a fallible player \( F^c \) in a lost btm position \( P \) of known depth, while White is infallible \( W^c \). Here, Depth to Conversion (DTC)\(^3\) is used (Tamplin, 2002) but DTM(ate), DTZ\(^4\) or DTR\(^5\) are equally valid. Below are listed some basic concepts and notation:

- \( P = \) a btm (wtm) position with \( n_a \) successors of depth \( a > 0 (\leq d-1) \), conversions to loss (win) being ignored
- \( d(P) = \) the depth of a position, taken to be \( \max(\text{depth}) + \omega \) for a drawn position, say \( \omega = 1 \)
- \( F^c = \) a fallible player with competence level \( c \): specific White and Black fallible players are \( FW^c \) and \( FB^c \)
- \( p_c(a) = \) the probability of loser \( F^c \) moving to a specific position of depth \( a = 0.a^c \), \( \theta \) ensuring \( \sum p_c(a) = 1 \)
- \( q_c(a) = \) the probability of winner \( F^c \) moving to a specific position of depth \( a = 0.a^c \)
- \( d_c(P) = \) \( E[\text{depth after move by loser } F^c \text{ from } P] = \sum a n_a p_c(a) a = (\sum a n_a a^c)(\sum a n_a a^c) \)
- \( e_c(P) = \) \( E[\text{depth after move by winner } F^c \text{ from } P] = \sum a n_a q_c(a) a = (\sum a n_a a^c)(\sum a n_a a^c) \)

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2 Scenario \( WFL \) has fallible White in a won wtm position against an infallible Black
3 \( DTC \) is expressed by a mate or by change of force, capture and/or promotion, on the board.
4 DTZ = Depth to move-count zeroing move, i.e. to mate, force-change or Pawn advance.
5 DTR = Depth by The Rule . which zeroes the move-count after a capture or advance of a Pawn (Haworth 2000, 2001).
3. FIRST RESULTS

The proofs involve only elementary algebra and analysis and are given in Appendix A.

**Theorem 1.** Scenario BFL: \(d_c(P) \leq d_{c+1}(P)\), i.e. increased competence increases expected depth. In this sense, \(FB_{c+1}\) is better than \(FB_c\). Equality occurs only if all successor positions have the same depth.

**Theorem 2.** Scenario BFL: \(d_c(P) \rightarrow d(P)\), i.e. \(\forall r, \exists c(r)\) such that \(c > c(r) \Rightarrow d_c(P) > d(P) - r\). e.g., for a losing defender with competence \(c \geq c(1)\), \(d_c(P) > d(P) - 1\).

**Theorem 3.** Scenario WFL: \(e_c(P) \geq e_{c+1}(P)\), i.e. increased competence lowers expected depth.

**Theorem 4.** Scenario WFL: \(e_c(P) \rightarrow d(P) - 1\), i.e. \(\forall r, \exists a_c(r)\) such that \(c > a_c(r) \Rightarrow d_c(P) < d(P) - 1 + r\).

**Theorem 5.** The numbers of positions \(n_a \geq 0\) may be modified by weights \(w \geq 0\). Weighting can reflect presumed search-depths or strategies, preferences and/or competencies on the part of the opponent (Jansen, 1992b, 1993), e.g. “plays only value preserving moves”, “will always identify a win or loss within 10 plies”, “tends to play moves with high tactical threat value” and so on.

4. BASIC USES OF FALLIBLE PLAYERS

The assumption of a fallible opponent enables a player to distinguish between metric-optimal\(^6\) moves. in the Browne-BELLE position BB2-22b, q.v. Table 1 which lists all positions used, the infallible defender Black has a choice between DTC-optimal Rf6 and Rf7. Figures 1-4 have White’s competence \(c\) as the x-axis and the expected depth, given \(c\), after White’s next move as the y-axis. Figure 1 shows that the expected depth after White’s reply to Rf7 is greater than after Rf6 – regardless of White’s actual competence \(c\).

<table>
<thead>
<tr>
<th>Position</th>
<th>Position</th>
</tr>
</thead>
<tbody>
<tr>
<td>BB1-01w</td>
<td>wKa8Qa5/bKf6Re8</td>
</tr>
<tr>
<td>BB2-01w</td>
<td>wKe8Qd8/bKe3Rc4</td>
</tr>
<tr>
<td>BB1-09b</td>
<td>wKd6Qd3/bKf4Re4</td>
</tr>
<tr>
<td>BB2-10b</td>
<td>wKe5Qg6/bKe3Rf4</td>
</tr>
<tr>
<td>BB2-12b</td>
<td>wKc5Qc3/bKe2Ra4</td>
</tr>
<tr>
<td>BB2-22b</td>
<td>wKe4Qe3/bKg4Rf5</td>
</tr>
<tr>
<td>BB2-24b</td>
<td>wKe4Qg3/bKf5Rf7</td>
</tr>
<tr>
<td>J1-1b</td>
<td>wKe3Qc5/bKf5Ra6</td>
</tr>
</tbody>
</table>

**Table 1.** KQKR positions.

At position BB2-10b, Black has a similar decision between DTC-optimals Ra4 and Re4. As Figure 2 shows, no move dominates across the whole range of \(c\): Ra4 is preferable against poor players but Re4 is significantly better for \(c > 6.3\). A probability distribution for \(c\) is therefore required: section 5 below addresses this issue.

The ranking of Black’s options for \(c = 0\) is determined by the average depth of White’s choices; as \(c \rightarrow \infty\) it is determined by the profile of White’s best responses to Black’s moves. Thus, it is not surprising that the two scenarios shown in Figures 1 and 2 can arise.

For position BB2-12b, Black has a choice between DTC-optimals Re4, Rf4, Rg4 or Rh4. The fallibility model shows that Rh4 is never the best option, while Rf4, Re4 and Rg4 are in turn best, respectively, for \(c \in [0, 19.7]\), \(c \in [19.7, 36]\) and \(c > 36\).

For position BB2-24b, Black has a choice between DTC-optimals Rb7, Rf6 and Rf8, q.v. Figure 3. The move Rf8 is just dominated throughout by Rf6 which, however, is best only for the range \(c \in (4.25, 28.5)\). Rb7 is best against both the zero-skill random mover and the near-infallible player.

\(^6\) ... the metric used here being DTC.
Finally, we recall Jansen’s (1993) suggestion that DTC-sub-optimal moves should also be considered: he illustrated this with another KQKR position, J1-1b. Certainly the move Kf6 which sacrifices one move in depth is at some stage better but, as Figure 4 shows, this is only for $c > 22.5$. Theorem 2 above helps limit the set of sub-optimal moves that need to be examined. For example, if it is known that $c > 22.5$, Kf6 must be preferred to any move ceding one or more moves in depth. Speculative play has greater potential in the harder endings when fallible opponents are likely to have less competence and lower apparent competence factors $c$.

## 5. IDENTIFYING COMPETENCE LEVELS

Perhaps the earliest demonstration of infallible player was in two KQKR games between the sporting volunteer GM Walter Browne and Ken Thompson’s BELLE (Fenner, 1979; Jansen, 1992a; Levy and Newborn, 1991), c.f. Appendix B. BELLE, defending, had Thompson’s newly calculated KQKR DTC EGT and apparently chose at random between equi-optimal moves.

The two initial positions were maxDTC positions with DTC = 31. BELLE claimed a 50-move draw in the first game but Browne, after a further bout of exhaustive preparation, won exactly on move 50 in the second game. It is intriguing to speculate as to whether BELLE could have squeezed out a second draw. Could it have set bigger problems for Browne by calculating which equi-optimal move to choose, or even by playing suboptimally (Jansen, 1993)? We address the first half of this question and speculate about the second.

Here we propose a way in which BELLE might have developed a probability distribution for Browne’s competence $c$, and the same technique can be used by an infallible observer when both players are fallible.
Let us suppose that BELLE had made a conventional, neutral, ‘know nothing’ initial assumption about Browne’s competence level $c$, for example that $c$ was equally likely to be 0, 10, 20, 30, 40 or 50. The initial $1/6$ probabilities could then be adjusted by the following rule of Bayesian inference:

$$\text{Posterior}_\text{Prob}(c) \propto \text{Prior}_\text{Prob}(c) \times \text{Prob}[\text{observed move} | c]$$

and then $\sum \text{Prob}(c). (\text{Expected depth} | c)$ could be calculated for each move option.

This computation, also reasonably assuming Browne would never choose a drawing or losing move, shows that his competence profiles in the two games are remarkably similar. The first fourteen moves elevate $E[c] = \sum \text{Prob}(c).c$ towards the maximum possible in the model before a few sub-optimal moves bring it sharply down.

However, in moves 20-22 and 32 of the second game, Browne progressed where he had stalled in game 1. Even so, the final $E[c]$ values for each game were similar, q.v. Figure 5. Browne ceded DTC depth as follows:

game 1: moves 6w (+1/1), 17w (+3/4), 18w (+2/6), 19w (+1/7), 20w (+1/8), 21w (+2/10), 22w (+1/11), 26w (+2/13), 31w (+1/14), 32w (+4/18), 33w (+2/20) ... 40w (+3/23), 41w (+2/25), 42w (+2/27)

game 2: moves 6w (+1/1), 16w (+1/2), 17w (+3/5), 19w (+3/8), 26w (+1/9), 27w (+2/11), 28w (+2/13), 30w (+1/14), 32w (+1/15), 33w (+1/16), 35w (+2/18), 44w (+1/19)

BELLE had equi-optimal choices in game 2 at moves 10b, 12b, 22b and 24b, hence their use in section 4 above. It picked the correct option only on move 22b. Assuming that Browne was roughly equivalent to player $F_{20}$, the opponent model shows a 0.3 move advantage for the best choices, so perhaps BELLE might have got a second draw. It is unlikely that speculative play would ever have been worthwhile.

In general, a set of fallible players, of which the set $\{F_c\}$ is one example, may be incorporated into a ‘PrOM’ Probabilistic Opponent-Modelling Strategy (Donkers, Uiterwijk and van den Herik, 2001).

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7 Relatively coarse-grain for convenience of calculation: a finer-grain initial $c$-profile is $c = 0(1)100$, probability 1/101.
6. EXPLOITING FALLIBLE PLAYERS

6.1 Endgame Practice Partners

The current 7-hour schedule of play is (40/2°, 20/1°, All/30'), involving fast third-phase rates of play to a finish. Good endgame technique is therefore increasingly important. The fallible players $F_c$ would make ideal practice partners, being available, tireless, house-trained, uncritical and non-deterministic in behaviour. They can comment on moves played, be set manually or automatically to a level of difficulty $c$, and even help you win if $c$ is set negative. These benefits recommend them as an adjunct to existing chess engines.

6.2 Changing the game result

The most significant use of the $F_c$ is perhaps to change the game score by differentiating between drawn positions. Attacker can pressure fallible defender and equally, defender can resist fallible attacker (Levy 1987, 1991; Nunn, 2002; Schaeffer 1991, 1997). Note that, as usual, the estimated values of positions may be backed up, in this case factored by the calculated probabilities of them occurring.

The suggestion that theoretically drawn positions should be given a depth of $\text{max}(\text{depth})+\omega$ was inevitably arbitrary on two counts. First, although $\omega$ must be positive to make draws less likely for an attacker than wins, it is not otherwise defined. Secondly, not all draws are the same. Blunder moves leading to capture and shallow 'draws in $n$' are less likely than play which merely retains the draw in the endgame.

6.3 Avoiding draw claims

Haworth (2000, 2001) establishes the DTR, Depth by The Rule, metric as necessary and sufficient for best avoiding $k$-move draw claims, whether $k = 50$ or not. If the opponent is identified as playing to the DTM or DTC metrics, it may be possible to induce them to sacrifice depth in DTR terms. However, in general, if the attacker assesses the opponent’s apparent competence as in section 6.2 above, this will inform their choice of equi-optimal or even sub-optimal moves.

6.4 Monte Carlo Simulation

From a position $P$, games $F_c-F_\infty$ may be played out to demonstrate how well $F_c$ can progress the win: similarly, games $F_\infty-F_c$ examine defensive capability. Engine-$F_c$ games can benchmark a chess engine’s endgame capability. Matches $F_c-F_c$ may be set up to determine whether, in general, an endgame is harder to win that to defend.

6.5 Assessing Endgame difficulty

Here, we attempt to derive some aggregate characteristics of an endgame as a whole rather than examining play from specific positions.

An endgame is modelled here as a Markov system: depth $i$ corresponds to state $i$, and endgame play is the Markov process. Averaged across the endgame, a White $F_c$ will move the endgame from depth $i$ to depth $j$ with probability $p_{ij,c}$, then $d_i = [p_{ij,c} | \text{fixed } i]$ is the depth probability vector after White’s move. Thus $MW_c = [p_{ij,c}]$ is a transition matrix for White’s moves: $MB_c$ is the corresponding matrix $[q_{ij,c}]$ for Black’s moves. For the infallible players, $MW_c$ has $p_{i,i-1,c} = 1$ and otherwise 0: $MB_c$ is the identity matrix. If $d_n$ is a depth probability vector after $n$ moves, $d_n(MW_c MB_c)^m$ is the same vector after a further $m$ moves by White $F_c$ and Black $F_b$.

The computation of the matrices $MW_c$ and $MB_c$ are feasible for 3-5-man endgames, given DTC depths $\leq 114$ and the the EGT sizes involved. With them, we may address questions such as:

- what is the expected depth $ED_{w,n}$ after $n$ moves of $F_c-F_\infty$ play from a given position?
- what is the expected number of moves $V_n$ required to win $F_c-F_\infty$ from a given position?

Two theorems, proved in Appendix A, and two likely conjectures follow.
Theorem 6. Let $E[d]$ = expected depth associated with probability vector $d$. Then $E[d,MW] \geq E[d,MW_{c+1}]$. Symmetrically, $E[d,MB] \leq E[d,MB_{c+1}]$. See Appendix A for the proof.

Theorem 7. If $u_i = \text{Prob[depth} = i\text{]} = 1$, then $E[u_i] = i$ and $\exists c^* \text{ s.t. } c > c^* \Rightarrow E[u_i,MW] < i < E[u_{i+1},MW] \forall i$.

Conjecture 1. There is a $w^*$ s.t. $w_2 > w_1 > w^* \Rightarrow ED_{w_1,n} \geq ED_{w_2,n}$ and $V_{w_1,n} \geq V_{w_2,n}$.

Conjecture 2. $\exists w = w_d(\text{endgame}) \text{ s.t. } w > w^* \Leftrightarrow F_w \text{ can secure all wins of depth } d$, i.e. $ED_{w,n} \rightarrow 0$.

Certainly, the graphs $ED_{w,n}$ from the aggregate maxDTC positions should indicate where the difficulties lie in depth terms. The parameters $\{w_d\}$ are characteristics of the endgame and may usefully indicate its relative difficulty.

7. SUMMARY

A number of questions arise as to how best to attack or defend against a fallible opponent. How best can a result be created or expedited? How competent is the opponent? This note has defined a reference spectrum of fallible endgame players whose use enables these and other questions to be addressed. Some reasonably tractable computations have been defined which could assess the aggregate difficulty of at least 3-5 man endgames.

Thanks goes to John Tamplin for many 3-5-man pawnless endgame DTC EGTs, including KQKR used here.

8. REFERENCES


APPENDIX A: PROOFS OF THEOREMS

Theorem 1. Scenario BFL: \( d(P) \leq d_{\ast,\ast}(P) \), i.e. expected depth after move increases with competence.

**Proof.** Let \( P \) be a position of depth \( d \), it is thought that \( d_{\ast,\ast}(P) \) is not constrained to be either integer or non-negative: \( F \), with negative skill are conceivable. Note also, in anticipation of Theorem 5, that the \( n_o \) are constrained only to be non-negative.

Theorem 2. Scenario BFL: \( d(P) \rightarrow d(P) \), i.e. for each \( r \), there is a \( c(r) \) such that \( d_{\ast}(P) > d(P) - r \).

**Proof.** Let \( d(P) = d : d(P) = (\Sigma n_o a^c \vdash \Sigma n_o a^c) \). Then \( E[\Sigma n_o a^c] = \Sigma n_o a^c/d(P) \) for \( a \leq d \).

Dividing numerator and denominator by \( d \)...

\[ d(P) = (\Sigma n_o a^c/d(P)) \]

\[ = (n_o d + \Sigma n_o a^c/d(P)) \]

\[ \rightarrow d = d(P) \text{ as } c \rightarrow \infty \]

Theorem 3. Scenario WFL: \( e(P) \geq e_{\ast,\ast}(P) \), i.e. expected depth after move decreases with competence.

**Proof.** Let \( e(P) = e : e(P) = (\Sigma n_o a^c \vdash \Sigma n_o a^c) \). Then \( E[\Sigma n_o a^c] = \Sigma n_o a^c/e(P) \) for \( a \geq e \).

Dividing numerator and denominator by \( e \)...

\[ e(P) = (\Sigma n_o a^c/e(P)) \]

\[ = (n_o e + \Sigma n_o a^c/e(P)) \]

\[ \rightarrow e = d(P) - 1 \text{ as } c \rightarrow \infty \]

Theorem 5. The numbers of positions \( n_o \geq 0 \) may be modified by weights \( w \geq 0 \).

**Proof.** The proofs of Theorems 1 & 3 require only that \( \forall a, b, n_o,n_a \geq 0 \). Thus, \( n_o \rightarrow w n_o \) is allowable.

Theorem 6. Let \( d \) be a depth probability vector and let \( E[d] = E[\Sigma c d_i u_i] = \text{expected depth associated with } d \).

Then \( E[d, MW] \geq E[d, MW_{\ast,\ast}] \). Symmetrically, \( E[d, MB] \leq E[d, MB_{\ast,\ast}] \).

**Proof.** This follows directly from Theorem 3 and, exchanging White for Black, from Theorem 1. \( E[d] = E[\Sigma c d_i u_i] \) is a linear function on the vector space \( \{d\} \).

Therefore we need only prove the theorem for a typical component of \( d \).

Let \( y_i = [0, ..., 1, 0, ..., 0] \), representing the fact that play is at any position of depth \( i \).

From theorem 3, we have \( e_i(P) \geq e_{\ast,\ast}(P) \).

Summing over all \( x \) positions of depth \( n \):

\[ \Sigma c e_i(P) \geq E[\Sigma c e_i(P)] / \Sigma (E[\Sigma c e_i(P)]) / \Sigma E[\Sigma c e_i(P)] \geq E[\Sigma c e_i(P)] \]

Therefore \( E[d, MW] \geq E[d, MW_{\ast,\ast}] \) for an arbitrary initial state vector \( d \).

Theorem 7. If \( y_i = \text{Prob}[\text{depth}=i]=1 \), then \( E[y_i] = i \) and \( \exists a c'-s.t. c + c' \Rightarrow E[y_i MW] < i < E[y_{i+1} MW] \forall i \).

**Proof.** Let \( P \) be a position of depth \( i \). By theorem 4, there is a \( c'(P) \) s.t. \( c + c' \Rightarrow e_i(P) < d(P) \).

Therefore, there is a \( c''(i) \) s.t. \( c + c'' \Rightarrow E[y_i MW] < i < E[y_{i+1} MW] \).

Therefore, there is a \( c'' \) s.t. \( c + c'' \Rightarrow E[y_i MW] < i < E[y_{i+1} MW] \forall i \).

It is thought that \( c'' \) is quite small.
APPENDIX B: THE BROWNE-BELLE GAMES

These games have been annotated with respect to the DTC metric as follows:

- `≡` only value-preserving move,
- `*' = only optimal move,
- `º` = only legal move,
- `{+/+i/j}` = move ceding `i` moves in depth, making a total loss of `j` moves in depth in the game


White ceded 27 moves in depth over moves 6, 17-22, 26, 31-33 and 40-42. Black had DTC-optimal choices on moves 8, 9, 12, 13, 20, 27 and 29 and, made the correct choice on only moves 8, 20 and 38.

Game 2: `{BB2-1w}: wKc8Qd8/bKc3Re4+w, dc = 31} 1. Kb7 Rb4+ 2. Ke6 Rc4º 3. Kb5 Rb4+º 4. Ka5 Re4º 5. Qd6 Rd4º 6. Qc5+ `{1+/2: Qd3} Ke6, Qc6 and Qh2 optimal} Kd3º 7. Kb5 Re4º 8. Qf6 Ke3º 9. Kc5 Re4º 10. Qg6 Ra4 `{Re4} 11. Qg3+ Ke2º 12. Qc3 Rf4 `{Re4, Rg4} 13. Kd5 Rh4º 14. Qe2+ Ke3º 15. Qd1+ Kf2º 16. Qd2+ `{1+/2: Qd3} Kf3º 17. Qe1+ `{3/+5: Qd3} Kg4º 18. Qd1+ Kf4º 19. Qe2+ `{3/+3: Qc1+} Kg4º 20. Qe4+ Kf5º 21. Ke4+ Kd6º 22. Ke4+ Kh5º 24. Qc3+ Kh5º 25. Ke4+ Kh5º 26. Ke6+ `{1+/9: Qd3} Rf6 27. Qa3+ `{2+/11: Qe5+} Rf8º 28. Qh3+ `{2+/13: Qg3} Kg5º 29. Qg3+ Kg4º 30. Qe5+ `{1+/14: Qe5+} Kg4º 31. Qh4º Kg5º 32. Ke5+ `{1+/15: Qd2+} Kd6º 33. Qh8+ `{1+/16: Qe2+ and Qh3 optimal} Kg5+º 34. Ke6º {and now, this pattern repeats at positions 37b, 40b, 43b and 46b} Kg4º 35. Qc8+ `{2+/18: Qc5} Kg5º 36. Qd7+ Kg5º 37. Ke5+ Kg4º 38. Kg7+ Kh4º 39. Kg6+ Kg4º 40. Ke4+ Rg2º 41. Qg6+ Kh3º 42. Qh5+ Kg3º 43. Ke3+ Rg1º 44. Qg5+ `{1+/19: Qe5} and now White cannot slip again} Kh2º 45. Qh4+ Kg2º 46. Ke2+ Ra1º 47. Qe4+ Kh3º 48. Qh7+ Kg3º 49. Kg7+ Kh3 50. Qx1º {Just in time! Bl. resigns} 1-0.

White ceded 19 moves in depth over moves 6, 16-17, 19, 26-28, 30, 32-33, 35 and 44. Black had DTC-optimal choices on moves 10, 12, 22 and 24, and made the correct choice on move 22 only.